

m68k

The chip itself

Motorola 68k

- The m68k is a 16/32-bit CISC microprocessor made by Motorola in 1979.
 - One of the last Big-Endian ISAs.
 - It was a popular Instruction Set in the 80s, powering Apple computers, Amigas, Sinclairs, X68000, Ataris, the Sega Genesis, SUN, and NeXT etc.
 - It went from 8MHz on the original 68000 to 133Mhz on the 68060.
 - Most of these had 24-bit or 32-bit addresses.
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M68k Assembly

- The m68k has 7 data registers, D0-D7, used to store 32-bits.
- It has 7 address registers, A0-A7 used to store addresses in memory.
- 1 register, Program Counter (PC), stores the current instruction being executed.
- 1 register, the Condition Code Register, (CCR) holds 5 bits to hold conditions about the last instruction which executed.
 - X (Extend)
 - N (Negative)
 - Z (Zero)
 - V (Overflow)
 - C (Carry)
- 3 main types of Ints:
 - (S)hort = 8-bits
 - (W)ord = 16-bits
 - (L)ong = 32-bits
- Addressing modes:
- Register: D0, A0.
- Address indirect: (A0); post-inc (A0)+; pre-dec -(A0).
- Displacement: (16,A0) adds signed 16-bit.
- Indexed: (8,A0,D1.w) or (0,A0,D1.l) (base + index*1 + disp).

- Absolute: (label).w (16-bit PC-space), (label).l (full 32-bit).
 - PC-relative: (16,PC), (8,PC,D1.w) (for PIC).
 - Immediate: #123, #-5, #0xDEAD.
-

M68k Assembly Continued

- Labels are used to refer to variables or functions:

```
.text
.even
.globl _start
_start:
    move.l #0, %d0    ; return code in D0
    rts              ; return
```

- Jumps include:
 - `jsr <subroutine>`: (J)ump (T)o (S)ubroutine - calls a function.
 - `bsr <subroutine>`: (B)ranch (T)o (S)ubroutine - pushes the next instruction to the stack and calls a function. `rts` is only after a branch.
 - `bra <subroutine>`: (B)Ranch (A)lways - jump forwards or backwards.
 - * `bra.s`: (B)Ranch (A)lways (S)hort - jump up to 126 bytes forwards/backwards.
 - * `bra.w`: (B)Ranch (A)lways (W)ord - jump up to 32766 bytes forwards/backwards.

Note that branch instructions are faster and smaller in memory, but limited in how far they can jump. `jsr` can jump anywhere.

- `move.(w/l) <src>, <dst>`: (Move) - move from `src` to `dst`.
- `movea.(w/l) <src>, AR`: (Move) (A)ddress - move `src` to address register.
- `cmp`: (C)o(m)p(arison) - sets flags after a comparison.
- `bcc`: lets you branch on a compare.
- `lea <src> <dst>`: (L)oad (E)ffective (A)ddress - Load a given address and place in `dst`.

```
lea    (4,A0), A1 ; add 4 to A0 and store address in A1
move.w (A1), D1 ; dereference A1 and put in D1
```

- Stack manipulation: stack grows downward:
 - Pushing: `move.w #18, -(sp)`
 - Popping: `addq.l #2, sp`
-

Running an m68k Machine

- There are two emulators I've used.
- Hatari
- AranyM
- These can run any kind of OS. I did not go with Linux (sorry) and ran AFROS, which is a TOS/Freemint compatible OS.



Freemint is a open source version of the MiNT operating system for Atari STs.

- Freemint's Github has the kernel freemint, two versions of libc (mintlib, libcmini), and other useful libs like gemlib

As well, there's a version of gcc with support for compiling to this target. Thorsten Otto's website has gcc compilers built to target the `m68k-atari-mint-gcc`. You can download these and start compiling C code to run on your emulator.

What about Rust?

- I tried to compile Rust to this target too: My custom target
- LLVM doesn't compile to this target well. There's some support for bare-metal (m68k-unknown-elf) but this doesn't give you any user space. It also outputs ELF files which don't run on our emulators out of the box anyway (they use a.out).
- Stack trace that shows the issue in LLVM's Instruction Selection:

```
/lib/librustc_driver-67dd58abd922ea00.so(+0x3c22f0f) [0x7faf1a822f0f]
/lib64/libc.so.6(+0x1a070) [0x7faf16a28070]
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(+0x763347d) [0x7faf1443347d]
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(+0x762d5fc) [0x7faf1442d5fc]
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(+0x7629bd2) [0x7faf14429bd2]
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(+0x7e9d192) [0x7faf14c9d192]
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(_ZN4llvm16SelectionDAGISel17CodeGenAndEmitDAGEv+0x8)
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(_ZN4llvm16SelectionDAGISel20SelectAllBasicBlocksERK16SelectionDAGISel)
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(_ZN4llvm16SelectionDAGISel20runOnMachineFunctionERK16SelectionDAGISel)
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(+0x8b7a6c9) [0x7faf1597a6c9]
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(_ZN4llvm13FPPassManager13runOnFunctionERNS_8FunctionPassManager)
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(_ZN4llvm13FPPassManager11runOnModuleERNS_6ModuleEv)
/lib/libLLVM.so.21.1-rust-1.93.0-nightly(_ZN4llvm6legacy15PassManagerImpl3runERNS_6ModuleEv)
```

```

/lib/librustc_driver-67dd58abd922ea00.so(+0x64b31c5) [0x7faf1d0b31c5]
/lib/librustc_driver-67dd58abd922ea00.so(+0x64b2dfb) [0x7faf1d0b2dfb]
/lib/librustc_driver-67dd58abd922ea00.so(+0x64b4ff5) [0x7faf1d0b4ff5]
/lib/librustc_driver-67dd58abd922ea00.so(_RINvNtNtCsaEKiI0nD17M_3std3sys9backtrace28__rust.
/lib/librustc_driver-67dd58abd922ea00.so(+0x6253539) [0x7faf1ce53539]
/lib/librustc_driver-67dd58abd922ea00.so(+0x6258def) [0x7faf1ce58def]
/lib64/libc.so.6(+0x71f54) [0x7faf16a7ff54]
/lib64/libc.so.6(+0xf532c) [0x7faf16b0332c]

```

Time to write C

- If you don't want to write any libc or anything, you can stop right here. Just use the C toolchain to compile to your emulator and you're good to go.
- I wanted to go further – I wanted to write my own version of libc to really understand the bits and bytes before main.

```

// How does our binary get argc and argv?
int main(int argc, char* argv[]) {
    printf("hello world\n");
    return 0; // what happens here?
}

```

Does anyone know?

start.ld

- First, we have to layout our binary.
- .text is where the instructions are for your program
- .data is where initialized global/static variables go. Don't edit them.
- .bss is where uninitialized global/static variables go.

```

OUTPUT_FORMAT(a.out-mintprg)
OUTPUT_ARCH(m68k)
ENTRY(_start)

```

```

SECTIONS {
    /* Start at 0 for PRG (BFD will add the PRG header). */
    . = 0;

    /* --- TEXT + RO --- */
    .text : ALIGN(2) {
        KEEP(*(.text._start))           /* ensure _start first */
        *(.text .text.*)
        *(.rodata .rodata.*)
    }
}

```

```

    KEEP(*(.init))
    KEEP(*(.fini))
}

/* --- DATA --- */
.data : ALIGN(2) {
    *(.data .data.*)
    *(.sdata .sdata.*)
    *(.init_array .init_array.*)
    *(.fini_array .fini_array.*)
}

/* --- BSS --- */
.bss (NOLOAD) : ALIGN(2) {
    *(.bss .bss.*)
    *(COMMON)
}

/* Strip sections not used on TOS */
/DISCARD/ : {
    *(.comment) *(.note*) *(.eh_frame*) *(.stab*) *(.debug*) *(.gnu.*)
}
}

```

crt0.S

- C's runtime, the stuff that runs before main.

```

.title "crt0.S for m68k-coff"
#define STACKSIZE 0x4000
/*
 * Define an empty environment.
 */
.data
.align 2
SYM (environ):
    .long 0
    .align 2
    .text
/*
 * These symbols are defined in C code, so they need to always be
 * named with SYM because of the difference between object file formats.
 */
    .extern SYM (main)
    .extern SYM (exit)

```

```

        .extern SYM (hardware_init_hook)
        .extern SYM (software_init_hook)
        .extern SYM (atexit)
        .extern SYM(__do_global_dtors)
/*
 * These values are set in the linker script, so they must be
 * explicitly named here without SYM.
 */
        .extern __stack
        .extern __bss_start
        .extern _end
/*
 * Set things up so the application will run. For historical reasons
 * this is called 'start'. We set things up to provide '_start'
 * as with other systems, but also provide a weak alias called
 * 'start' for compatibility with existing linker scripts.
 */
        .global SYM (start)
        .weak SYM (start)
        .set SYM (start),SYM(_start)
        .global SYM (_start)

```

crt0.S continued

```

SYM (_start):
    /* See if user supplied their own stack (__stack != 0). If not, then
     * default to using the value of %sp as set by the ROM monitor.
     */
    movel    IMM(__stack), a0
    cmpl     IMM(0), a0
    jbeq     1f
    movel    a0, sp
1:
    /* set up initial stack frame */
    link     a6, IMM(-8)
/*
 * zero out the bss section.
 */
    movel    IMM(__bss_start), d1
    movel    IMM(_end), d0
    cmpl     d0, d1
    jbeq     3f
    movl     d1, a0
    subl     d1, d0

```

```

    subq1    IMM(1), d0
2:
    clrb     (a0)+
    subq1    IMM(1), d0
    jbp1     2b

/*
 * call the main routine from the application to get it going.
 * main (argc, argv, environ)
 * we pass argv as a pointer to NULL.
 */
    movel    IMM (__FINI_SECTION__), (sp)
    PICCALL  SYM (atexit)
    PICCALL  __INIT_SECTION__
    pea      0
    PICPEA   SYM (environ), a0
    pea      sp@4
    pea      0
    PICCALL  SYM (main)
    movel    d0, sp@-

```

Now we're in main

That leaves one question: How does writing to the console work?

```

int main(int argc, char* argv[]) {
    if (argc > 1) {
        puts(argv[1]); // but how does this work?
    }
}

```

- We need something called a System Call.
-

System Calls

- System calls ask the kernel to do something on the user's behalf.
- Operating systems provide 3 things: Virtualization, Concurrency, Persistence.
- We cannot have any of these if users can do what they want to.
- Thus there's a need for a contract between the user and the OS.

Freemint's System Call for Fwrite

We can call this special `trap #1` after loading our arguments onto the stack and the OS will write for us.

```

pea      buf           ; Offset 8
move.l   count,-(sp)    ; Offset 4
move.w   handle,-(sp)   ; Offset 2
move.w   #64,-(sp)      ; Offset 0
trap     #1             ; GEMDOS
lea      $C(sp),sp      ; Correct stack

```

Translated to C:

```

/* Fwrite - Write bytes to a file handle. */
static inline int32_t Fwrite(int16_t handle, int32_t count, const void *buf) {
    return (int32_t)trap_1_wll(0x40, handle, (long)count, (long)buf);
}

```

System Calls Continued

```

static inline long trap_1_wll(short n, short a, long b, long c) {
    register long retval __asm__("d0");
    short _a = (short)(a);
    long _b = (long)(b);
    long _c = (long)(c);

    __asm__ volatile("movl    %4,%%sp@-\n\t"
                     "movl    %3,%%sp@-\n\t"
                     "movw    %2,%%sp@-\n\t"
                     "movw    %1,%%sp@-\n\t"
                     "trap     #1\n\t"
                     "lea %%sp@(12),%%sp"
                     : "=r"(retval) /* outputs */
                     : "g"(n), "r"(_a), "r"(_b), "r"(_c) /* inputs */
                     : "d1", "d2", "a0", "a1", "a2", "cc", "memory" /* clobbered */
                     );
    return retval;
}

```

Implementing fwrite

This code:

```

size_t fwrite(const void *ptr, size_t size, size_t nmemb, FILE *stream) {
    if (size == 0 || nmemb == 0)
        return 0;

    size_t total_bytes = size * nmemb;
    long ret = Fwrite(stream->handle, total_bytes, ptr);
    if (ret <= 0)

```

```

    return 0;

    return ret / size;
}

```

Turns into this:

```
$ m68k-atari-mint-objdump -dr build/objs/stdio/fwrite.o
```

```
build/objs/stdio/fwrite.o:      file format a.out-zero-big
```

Disassembly of section .text:

```

00000000 <_fwrite>:
   0:  48e7 3c20      moveml %d2-%d5/%a2,%sp@-
   4:  262f 001c      movel %sp@(28),%d3
   8:  202f 0020      movel %sp@(32),%d0
  c:  4a83          tstl %d3
  e:  6608          bnes 18 <.L2>

00000018 <.L2>:
 18:  4a80          tstl %d0
 1a:  67f4          beqs 10 <.L4>
 1c:  206f 0024      moveal %sp@(36),%a0
 20:  2810          movel %a0@,%d4
 22:  2f00          movel %d0,%sp@-
 24:  2f03          movel %d3,%sp@-
 26:  4eb9 0000 0000  jsr 0 <_fwrite>
 28:  32           __mulsi3      # what the
 2c:  508f          addql #8,%sp
 2e:  2a2f 0018      movel %sp@(24),%d5 # Here's the syscall
 32:  2f05          movel %d5,%sp@-   # We load the buffer
 34:  2f00          movel %d0,%sp@-   # Then the length
 36:  3f04          movew %d4,%sp@-   # then the handle
 38:  3f3c 0040      movew #64,%sp@-   # Then #64, syscall for fwrite
 3c:  4e41          trap #1          # trap
 3e:  4fef 000c      lea %sp@(12),%sp  # correct the stack
...

```

What is __mulsi3?

```
$ m68k-atari-mint-objdump -D /usr/lib64/gcc/m68k-atari-mint/12/libgcc.a | rg __mulsi3 -A 20
00000000 <__mulsi3>:
   0:  302f 0004      movew %sp@(4),%d0
   4:  c0ef 000a      muluw %sp@(10),%d0

```

```

8:  322f 0006      movew %sp@(6),%d1
c:  c2ef 0008      muluw %sp@(8),%d1
10: d041          addw %d1,%d0
12: 4840          swap %d0
14: 4240          clrw %d0
16: 322f 0006      movew %sp@(6),%d1
1a: c2ef 000a      muluw %sp@(10),%d1
1e: d081          addl %d1,%d0
20: 4e75          rts

```

- The m68k doesn't have hardware instructions for multiplying two 32-bit numbers (`mul.l` doesn't exist) on the original m68k – it was added to the 68020.
- Since I'm targeting the 68k, it's added in for me.

Let's implement some calls

How would you implement `memcpy`, which copies `n` bytes from `src` to `dest`?

```
void *memcpy(void *restrict dest, const void *restrict src, size_t n);
```

A simple implementation

```

void *memcpy(void *restrict dest, const void *restrict src, size_t n) {
    for (; n; n--) *d++ = *s++;
    return dest;
}

```

This is Musl's version

```

void *memcpy(void *restrict dest, const void *restrict src, size_t n)
{
    unsigned char *d = dest;
    const unsigned char *s = src;

#ifdef __GNUC__

    if __BYTE_ORDER == __LITTLE_ENDIAN
#define LS >>
#define RS <<
    else
#define LS <<
#define RS >>

```

```

#endif

typedef uint32_t __attribute__((__may_alias__)) u32;
uint32_t w, x;

for (; (uintptr_t)s % 4 && n; n--) *d++ = *s++;

if ((uintptr_t)d % 4 == 0) {
    for (; n>=16; s+=16, d+=16, n-=16) {
        *(u32*)(d+0) = *(u32*)(s+0);
        *(u32*)(d+4) = *(u32*)(s+4);
        *(u32*)(d+8) = *(u32*)(s+8);
        *(u32*)(d+12) = *(u32*)(s+12);
    }
    if (n&8) {
        *(u32*)(d+0) = *(u32*)(s+0);
        *(u32*)(d+4) = *(u32*)(s+4);
        d += 8; s += 8;
    }
    if (n&4) {
        *(u32*)(d+0) = *(u32*)(s+0);
        d += 4; s += 4;
    }
    if (n&2) {
        *d++ = *s++; *d++ = *s++;
    }
    if (n&1) {
        *d = *s;
    }
    return dest;
}

```

Continued

```

if (n >= 32) switch ((uintptr_t)d % 4) {
case 1:
    w = *(u32 *)s;
    *d++ = *s++;
    *d++ = *s++;
    *d++ = *s++;
    n -= 3;
    for (; n>=17; s+=16, d+=16, n-=16) {
        x = *(u32*)(s+1);
        *(u32*)(d+0) = (w LS 24) | (x RS 8);
    }
}

```

```

        w = *(u32 *)(s+5);
        *(u32 *)(d+4) = (x LS 24) | (w RS 8);
        x = *(u32 *)(s+9);
        *(u32 *)(d+8) = (w LS 24) | (x RS 8);
        w = *(u32 *)(s+13);
        *(u32 *)(d+12) = (x LS 24) | (w RS 8);
    }
    break;
case 2:
    w = *(u32 *)s;
    *d++ = *s++;
    *d++ = *s++;
    n -= 2;
    for (; n>=18; s+=16, d+=16, n-=16) {
        x = *(u32 *)(s+2);
        *(u32 *)(d+0) = (w LS 16) | (x RS 16);
        w = *(u32 *)(s+6);
        *(u32 *)(d+4) = (x LS 16) | (w RS 16);
        x = *(u32 *)(s+10);
        *(u32 *)(d+8) = (w LS 16) | (x RS 16);
        w = *(u32 *)(s+14);
        *(u32 *)(d+12) = (x LS 16) | (w RS 16);
    }
    break;
case 3:
    w = *(u32 *)s;
    *d++ = *s++;
    n -= 1;
    for (; n>=19; s+=16, d+=16, n-=16) {
        x = *(u32 *)(s+3);
        *(u32 *)(d+0) = (w LS 8) | (x RS 24);
        w = *(u32 *)(s+7);
        *(u32 *)(d+4) = (x LS 8) | (w RS 24);
        x = *(u32 *)(s+11);
        *(u32 *)(d+8) = (w LS 8) | (x RS 24);
        w = *(u32 *)(s+15);
        *(u32 *)(d+12) = (x LS 8) | (w RS 24);
    }
    break;
}
if (n&16) {
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
}

```

```

if (n&8) {
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
}
if (n&4) {
    *d++ = *s++; *d++ = *s++; *d++ = *s++; *d++ = *s++;
}
if (n&2) {
    *d++ = *s++; *d++ = *s++;
}
if (n&1) {
    *d = *s;
}
return dest;
#endif

for (; n; n--) *d++ = *s++;
return dest;
}

```

What did we learn?

- Maybe Intel shouldn't have won
- Retrocomputing can be fun
- Don't write your own libc, it'll take me months if I continue on